

SSL für alle

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How to get good SSL security



Agenda

What is HTTPS / SSL

Attacks on SSL

Attacks on related technology

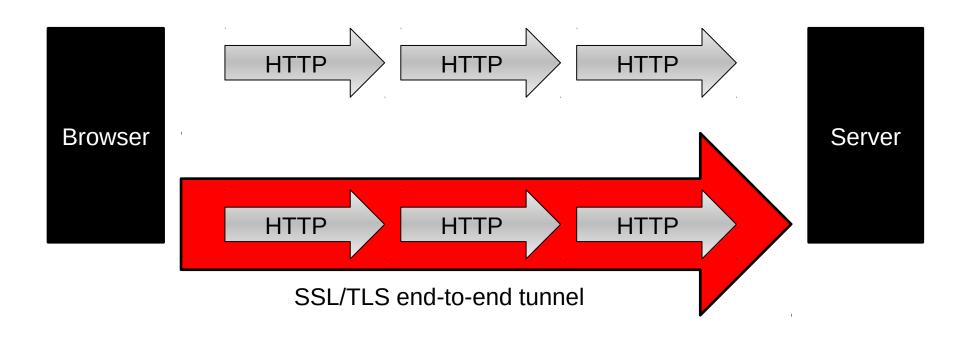
Coding SSL/TLS

What is HTTPS and SSL/TLS?





What is HTTPS?



What does SSL/TLS offer?

- Confidentiality
- Integrity / tamper resistance
- Authentication
 - -Server authentication

- Authenticated handshake.
 Communication is only MACed.
 DOES NOT offer non
 repudiation
- -Client authentication (option rarely used)
- Cryptographic agility
 - -Negotiate mutually accepted cipher suits



SSL/TLS versions

- SSL 3.0 (Netscape)
 - -1996, first widely used SSL. 1.0 & 2.0 broken.
- TLS 1.0
 - -1999, IETF standardization, & security fixes.
- TLS 1.1
 - –2006, CBC Cipher Block Chaining improvements
- TLS 1.2
 - -2008, Authenticated Encryption suites, AES, and many other security improvements.



SSL/TLS not perfect

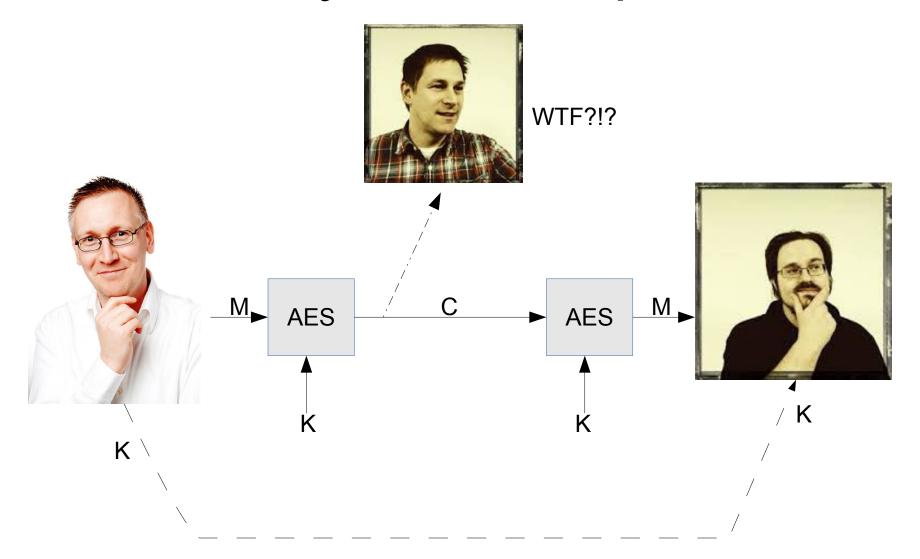
- SSL/TLS is not particularly well designed
 - -Evolution, with many hotfixes and workarounds
 - -Compatibility and legacy
 - Horribly slow adoption rate of new protocol versions and more secure cipher suites
- MAC is performed on plaintext not ciphertext
 - Most suites are by design vulnerable to chosen ciphertext attacks





Algorithms and Cipher Suites





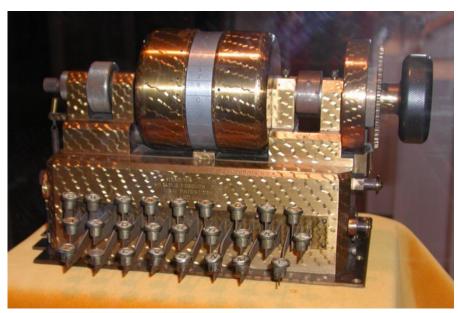


- Bulk encryption
 - AES, DES/3DES, Camellia, RC2, RC4
 - Fast, efficient in SW and HW
- Provides confidentiality NOT integrity
 - Encrypted message can be changed
- Must transfer secret key to receiver
 - How to protect the secret key during transfer?
 - Classic cipher with classic problem





Jefferson (1795)



Hebern (1918)





Hagelin C-36



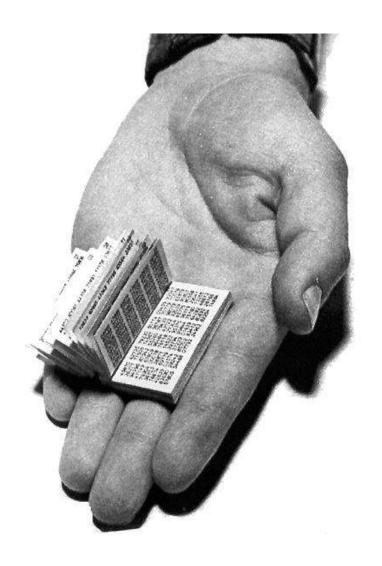
Enigma



Fialka

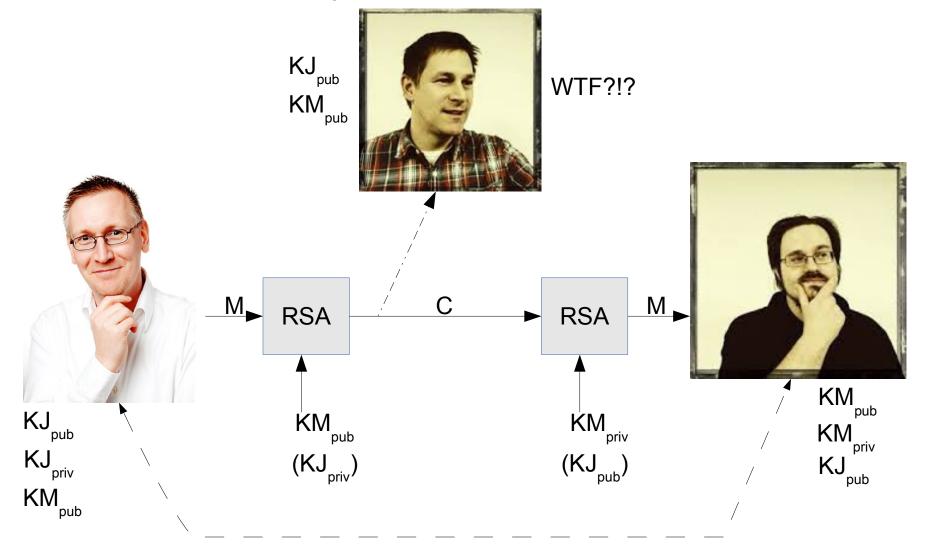


Keys for symmetric ciphers











- Public key encryption session init
 - RSA, ECC, El Gamal
 - 1000x-10000x slower than symmetric encryption
 - Complex math hard for embedded systems
- Provides confidentiality
 - Can provide integrity, origin authentication
 - Can provide exchange of secret keys (D-H)
- Must transfer public key to sender
 - No need to protect the key but need to trust the key
 - CAs provides trust by proxy (assumed trust in CA)



Hash functions



0x557e8e7ed1534946462f4136623947ca

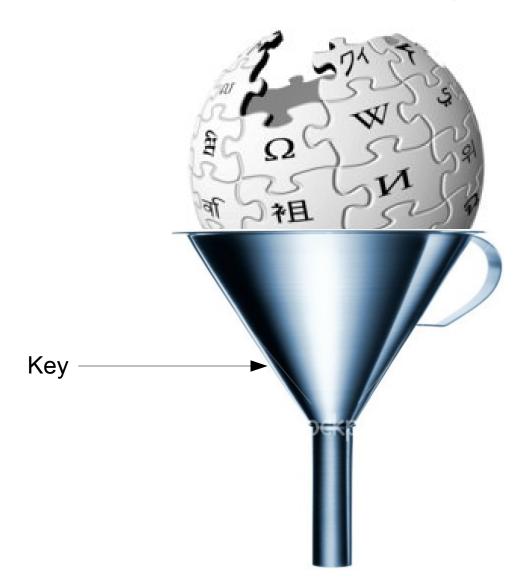


Hash functions

- Variable data size in, fixed size data out
 - Fingerprint, digest, hash related to input
- Keyless function
 - Security is based on collision resistance
- Provides data integrity
 - Detect presence of changes (errors) in the data



MACs



0x557e8e7ed1534946462f4136623947ca



MACs

Message Authentication Codes

- Variable data size in, fixed size data out
 - HMAC, OMAC, UMAC, CBC-MAC
 - Built using hash functions, block ciphers etc.
 - Fingerprint, digest, hash related to input
- Keyed function
 - Security is based on the secrecy of the key
 - Must transfer key to recipient
- Provides data integrity and authentication
 - Detect presence of changes (errors) in the data
 - Validate that the data is from the owner of the key



Cipher suites in my OpenSSL

```
js@secworks82.gotanet.se:/Users/js>openssl ciphers -v 'ALL:COMPLEMENTOFALL'
ADH-SEED-SHA
                                            Au=None Enc=SEED(128) Mac=SHA1
                         SSLv3 Kx=DH
DHE-RSA-SEED-SHA
                         SSLv3 Kx=DH
                                            Au=RSA
                                                    Enc=SEED(128) Mac=SHA1
DHE-DSS-SEED-SHA
                         SSLv3 Kx=DH
                                                    Enc=SEED(128) Mac=SHA1
                                            Au=DSS
SEED-SHA
                                            Au=RSA
                                                    Enc=SEED(128) Mac=SHA1
                         SSLv3 Kx=RSA
ADH-AES256-SHA
                         SSLv3 Kx=DH
                                            Au=None Enc=AES(256)
                                                                   Mac=SHA1
DHE-RSA-AES256-SHA
                                            Au=RSA
                                                    Enc=AES(256)
                         SSLv3 Kx=DH
                                                                   Mac=SHA1
DHE-DSS-AES256-SHA
                         SSLv3 Kx=DH
                                            Au=DSS
                                                    Enc=AES(256)
                                                                   Mac=SHA1
AES256-SHA
                                                    Enc=AES(256)
                         SSLv3 Kx=RSA
                                            Au=RSA
                                                                   Mac=SHA1
ADH-AES128-SHA
                                            Au=None Enc=AES(128)
                         SSLv3 Kx=DH
                                                                   Mac=SHA1
DHE-RSA-AES128-SHA
                                                    Enc=AES(128)
                         SSLv3 Kx=DH
                                            Au=RSA
                                                                   Mac=SHA1
                         SSLv3 Kx=DH
                                                    Enc=AES(128)
DHE-DSS-AES128-SHA
                                            Au=DSS
                                                                   Mac=SHA1
AES128-SHA
                         SSLv3 Kx=RSA
                                            Au=RSA
                                                    Enc=AES(128)
                                                                   Mac=SHA1
ADH-DES-CBC3-SHA
                         SSLv3 Kx=DH
                                            Au=None Enc=3DES(168) Mac=SHA1
ADH-DES-CBC-SHA
                                            Au=None Enc=DES(56)
                         SSLv3 Kx=DH
                                                                   Mac=SHA1
EXP-ADH-DES-CBC-SHA
                         SSLv3 Kx=DH(512)
                                            Au=None Enc=DES(40)
                                                                   Mac=SHA1 export
                         SSLv3 Kx=DH
ADH-RC4-MD5
                                            Au=None Enc=RC4(128)
                                                                   Mac=MD5
EXP-ADH-RC4-MD5
                         SSLv3 Kx=DH(512)
                                            Au=None Enc=RC4(40)
                                                                   Mac=MD5
                                                                            export
EDH-RSA-DES-CBC3-SHA
                                            Au=RSA
                                                    Enc=3DES(168) Mac=SHA1
                         SSLv3 Kx=DH
                                                                   Mac=SHA1
EDH-RSA-DES-CBC-SHA
                         SSLv3 Kx=DH
                                            Au=RSA
                                                    Enc=DES(56)
                                                                   Mac=SHA1 export
EXP-EDH-RSA-DES-CBC-SHA SSLv3 Kx=DH(512)
                                            Au=RSA
                                                    Enc=DES(40)
EDH-DSS-DES-CBC3-SHA
                                                    Enc=3DES(168) Mac=SHA1
                         SSLv3 Kx=DH
                                            Au=DSS
EDH-DSS-DES-CBC-SHA
                                            Au=DSS
                                                    Enc=DES(56)
                         SSLv3 Kx=DH
                                                                   Mac=SHA1
EXP-EDH-DSS-DES-CBC-SHA SSLv3 Kx=DH(512)
                                            Au=DSS
                                                    Enc=DES(40)
                                                                   Mac=SHA1 export
DES-CBC3-SHA
                         SSLv3 Kx=RSA
                                            Au=RSA
                                                    Enc=3DES(168) Mac=SHA1
DES-CBC-SHA
                         SSLv3 Kx=RSA
                                            Au=RSA
                                                    Enc=DES(56)
                                                                   Mac=SHA1
EXP-DES-CBC-SHA
                         SSLv3 Kx=RSA(512)
                                            Au=RSA
                                                    Enc=DES(40)
                                                                   Mac=SHA1 export
EXP-RC2-CBC-MD5
                         SSLv3 Kx=RSA(512) Au=RSA
                                                    Enc=RC2(40)
                                                                   Mac=MD5
                                                                            export
RC4-SHA
                         SSLv3 Kx=RSA
                                            Au=RSA
                                                    Enc=RC4(128)
                                                                   Mac=SHA1
RC4-MD5
                         SSLv3 Kx=RSA
                                            Au=RSA
                                                    Enc=RC4(128)
                                                                   Mac=MD5
EXP-RC4-MD5
                         SSLv3 Kx=RSA(512)
                                            Au=RSA
                                                    Enc=RC4(40)
                                                                   Mac=MD5
                                                                            export
DES-CBC3-MD5
                         SSLv2 Kx=RSA
                                            Au=RSA
                                                    Enc=3DES(168) Mac=MD5
DES-CBC-MD5
                                                    Enc=DES(56)
                         SSLv2 Kx=RSA
                                            Au=RSA
                                                                   Mac=MD5
EXP-RC2-CBC-MD5
                         SSLv2 Kx=RSA(512) Au=RSA
                                                    Enc=RC2(40)
                                                                   Mac=MD5
                                                                            export
RC2-CBC-MD5
                         SSLv2 Kx=RSA
                                                    Enc=RC2(128)
                                            Au=RSA
                                                                   Mac=MD5
EXP-RC4-MD5
                         SSLv2 Kx=RSA(512)
                                            Au=RSA
                                                    Enc=RC4(40)
                                                                   Mac=MD5
                                                                            export
                                                    Enc-BC4(128)
RC4-MD5
                         SSLv2 Kx=RSA
                                            Au=RSA
                                                                   Mac=MD5
                                                    Enc=None
NULL-SHA
                         SSLv3 Kx=RSA
                                                                   Mac=SHA1
                                            Au=RS/
NULL-MD5
                         SSLv3 Kx=RSA
                                            Au=RSN
                                                    Enc=None
                                                                   Mac=MD5
```

NULL = NO cipher!





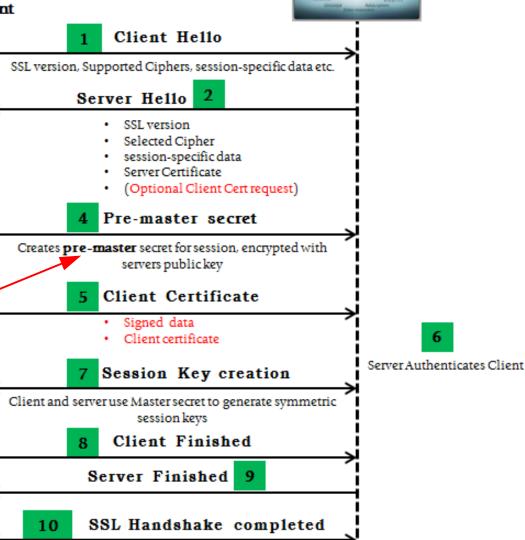


The SSL/TLS handshake

3

Client authenticates server

The session key!



Exchange Messages (Encrypted)

Source: Internet



Session keys in SSL/TLS

Client generates session master secret

 Client send secret to server. The secret is protected using the servers public key

 Session keys are then derived by client and server (encryption, MAC)

If the private key of the server is lost all previous sessions can be decrypted by extracting their master secrets



Perfect Forward Secrecy

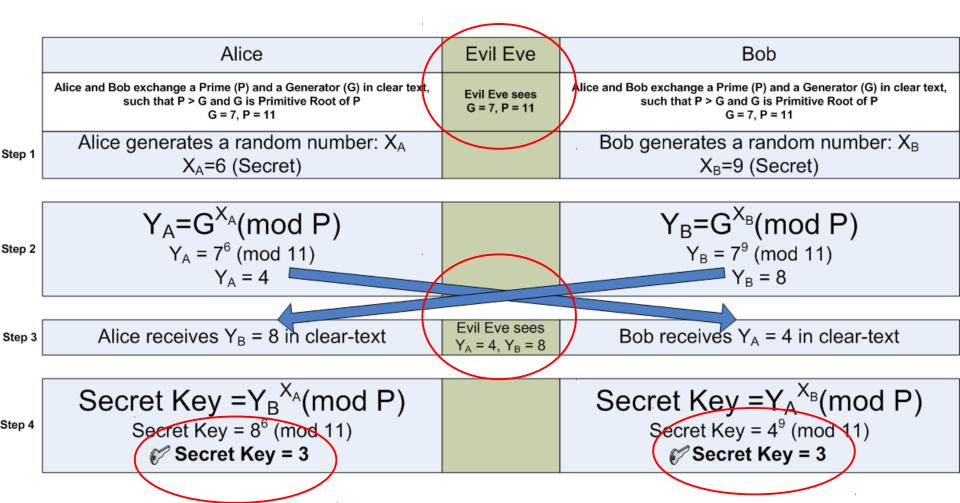
(Or simply Forward Secrecy)

- No master secret is transferred to server
 - Client and server agrees on common secret
 - Communicates using public messages
 - Diffie-Hellman key exchange
- After the session all secrets are discarded

Even if the private key of the server is lost all previous sessions are protected



Diffie Hellman Key Exchange



Copyright ©2005, Saqib Ali



Forward Secrecy at Twitter

Friday, November 22, 2013 | By Jacob Hoffman-Andrews (@j4cob) [21:39 UTC]

Tweet

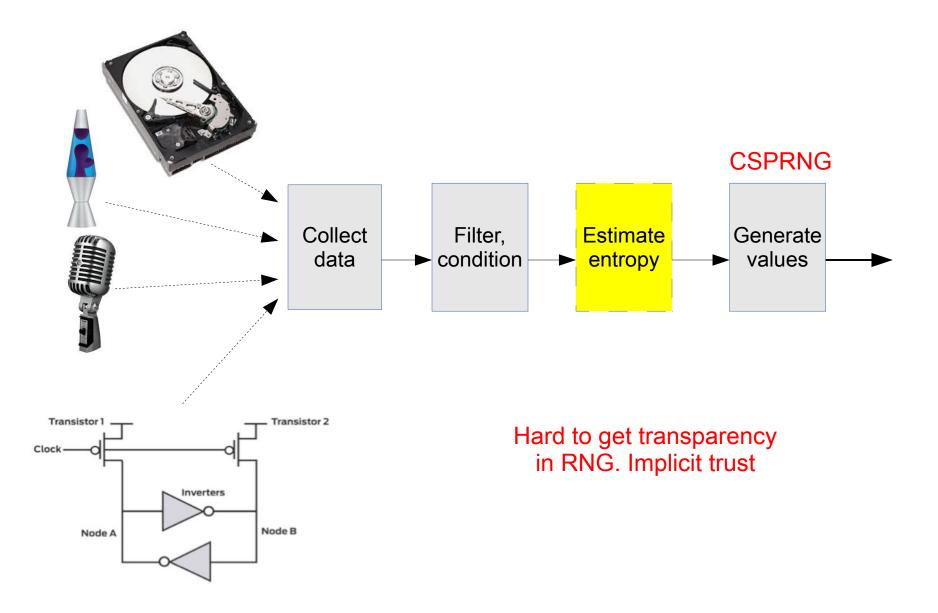
As part of our continuing effort to keep our users' information as secure as possible, we're happy to announce that we recently enabled forward secrecy for traffic on twitter.com, api.twitter.com, and mobile.twitter.com. On top of the usual confidentiality and integrity properties of HTTPS, forward secrecy adds a new property. If an adversary is currently recording all Twitter users' encrypted traffic, and they later crack or steal Twitter's private keys, they should not be able to use those keys to decrypt the recorded traffic. As the Electronic Frontier Foundation points out, this type of protection is increasingly important on today's Internet.



It is all about the keys!



Random Number Generation

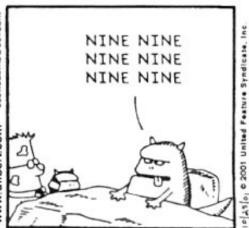


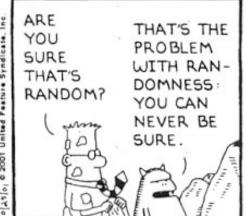


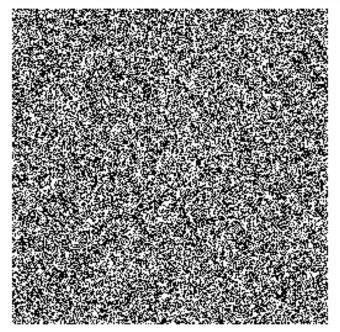
The problem with RNGs

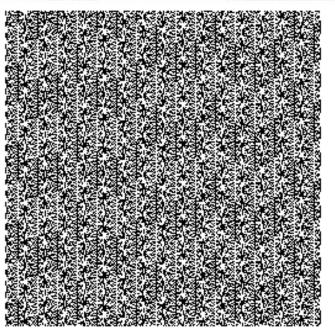
DILBERT By Scott Adams











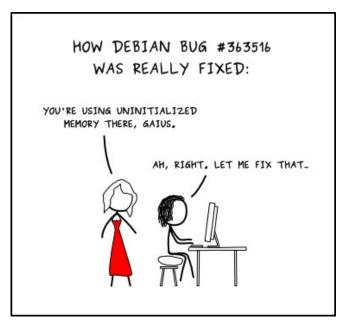


Debian key generator

Linux Debian (Etch, Lenny, Sid) 2006-2008

```
MD_Update(&m buf, j);
[ .. ]
MD_Update(&m buf, j); /* purify complains */
```

Security audit in Debian using Valgring, Purify OKed by OpenSSL dev

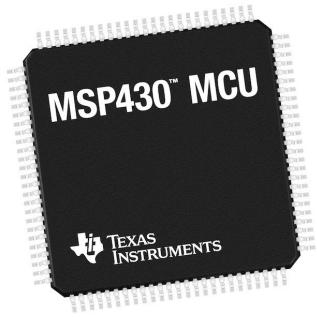


No mixing of random values during init

Process ID became seed At most ~17 bits strength

Affected SSL/TLS and a lot of applications





Z-Stack

air% hexdumptrandom bin | greps-fcolor "7c e1 e8 4e f4 87"
00000000 02 01 00 60 e8 2e 7c e1 e8 4e f4 87 62 49 56 fe
0008000 01 00 60 e8 2e 7c e1 e8 4e f4 87 62 49 56 fe 80
0010000 00 60 e8 2e 7c e1 e8 4e f4 87 62 49 56 fe 80 00
0018000 60 e8 2e 7c e1 e8 4e f4 87 62 49 56 fe 80 00 60
air% Source: T Goodspeed



Key length and strength

Symmetric	RSA/DLOG	EC
48	480	96
50	512	100
56	640	112
62	768	124
64	816	128
73	1024	146
80	1248	160
89	1536	178
103	2048	206
112	2432	224
128	3248	256
141	4096	282
160	5312	320
192	7936	384
256	15424	512

Source: Ecrypt II Yearly Report (2012) http://www.ecrypt.eu.org/documents/D.SPA.20.pdf

	OWASP	
	Open Web Application Security Project	

Table 7.4: Security levels (symmetric equivalent).

UMHSP		rable 7.4. Security levels (symmetric equivalent).					
Open Web Application Security Project	Security	Security	Protection	Comment			
	Level	(bits)					
	1.	32	Attacks in "real-time"	Only acceptable for			
			by individuals	auth. tag size			
	2.	64	Very short-term	Should not be used for			
			protection against	confidentiality in new			
			small organizations	systems			
	3.	72	Short-term protection	•			
4.			against medium				
			organizations, medium-				
			term protection against				
			small organizations				
	4.	80		n Smallest general-purpose			
			against agencies, long-	level, ≤ 4 years protection			
			term prot. against small	(E.g. use of 2-key 3DES,			
			organizations	< 2 ⁴⁰ plaintext/ciphertexts)			
	5 .	96	Legacy standard level	2-key 3DES restricted			
DCA 0040				to $\sim 10^6$ plaintext/ciphertexts,			
RSA 2048 -				≈ 10 years protection			
	6.	112	Medium-term protection	≈ 20 years protection			
			-	(E.g. 3-key 3DES)			
	7.	128	Long-term protection	Good, generic application-			
				indep. recommendation,			
				$\approx 30 \text{ years}$			
	8.	256	"Foreseeable future"	Good protection against			
				quantum computers unless			
ce: Ecrypt II Yearly F				Shor's algorithm applies.			
//www.ecrvpt.eu.org/	documents/E).SPA.20.pd	f	0 11			

Source: Ecrypt II Yearly R http://www.ecrypt.eu.org/documents/D.SPA.20.pdf



Table 7.1: Minimum symmetric key-size in bits for various attackers.

Attacker	Budget	Hardware	Min security
"Hacker"	0	PC	58
	< \$400	PC(s)/FPGA	63
	0	"Malware"	77
Small organization	\$10k	PC(s)/FPGA	69
Medium organization	300k	FPGA/ASIC	69
Large organization	\$10M	FPGA/ASIC	78
Intelligence agency	\$300M	ASIC	84



Side note: Export rules

- Waasenaar agreement
- Limits the usage of strong encryption
 - -Symmetric keys: 56 bits
 - -RSA keys: 512 bits
 - -Elliptic Curve keys: 112 bits
- Anything above these limits requires registration or permit
 - -EU, USA, Japan etc registration
 - -North Korea, Iran, China permit (or NO)

Also depends on usage End users vs components, equipment, market etc



Trust Stores

- Where does your app/OS find the CA cert?
- Trust Stores the root of cert validation
 - DB of CA certs in system, browser, libs
 - Mozilla
 - Apple OSX, iOS
 - Microsoft
 - Google Chrome, Oracle/Java
- Can you trust it?
 - Not very transparent, easy to check
 - Source code, blobs, Excel files
 - No info when the stores are updated and why
 - NSA root?

We try to observe the trust stores:

https://github.com/kirei/catt



Trust Store examples

Local CAs

CertiSign	Brazil	Certisign - Autoridade Certificadora - AC2	1024	MD5	Tuesday, June 26, 2018 5:00:00 PM

Government CAs

China Internet Network Information Center (CNNIC)	China	CNNIC Root	2048	SHA1	4/15/2027
China Internet Network Information Center (CNNIC)	China	China Internet Network Information Center EV	2048	SHA1	Saturday, August 31, 2030 12:11:25 AM
		_			

Big boys being bad

VeriSign	USA	VeriSign Class 1 Public Primary Certification Authority	1024	CUA1	Wednesday, August 02, 2028 4:59:59 PM



Agenda

What is HTTPS / SSL

Attacks on SSL

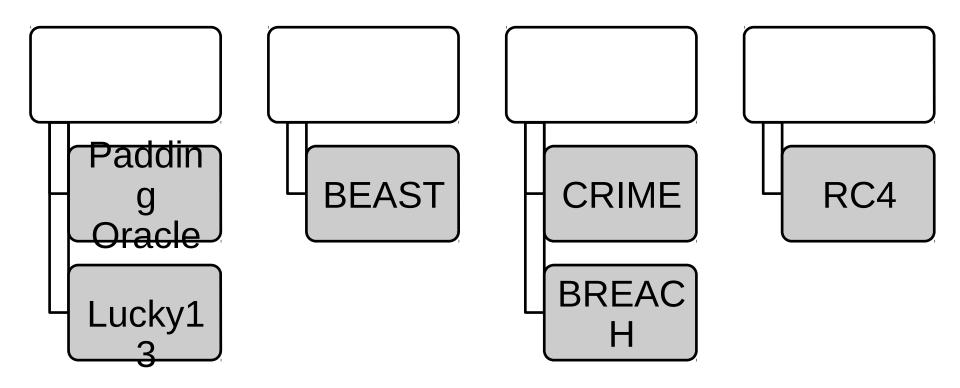
Attacks on related technology

Coding SSL/TLS

Cryptanalytic attacks against SSL

A. Shamir: "Cryptography is typically bypassed, not penetrated"







Padding Oracle & Lucky13

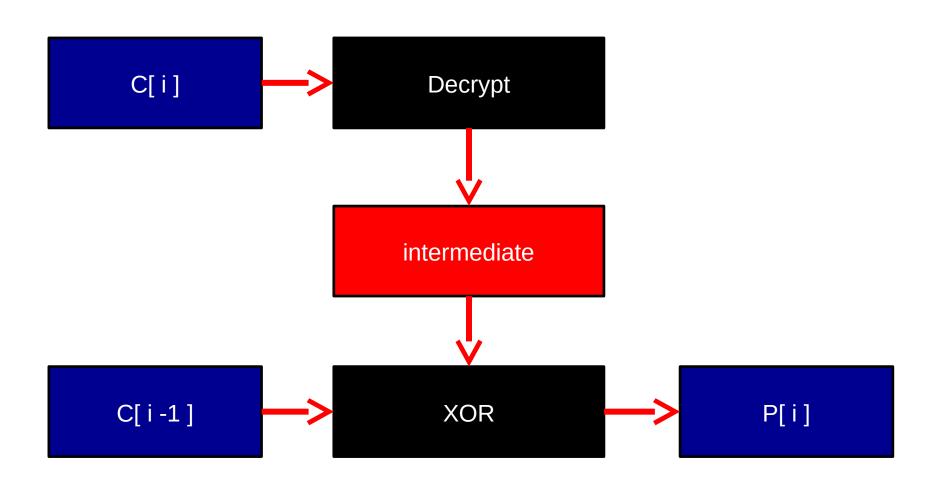
Cryptanalytic attacks against SSL



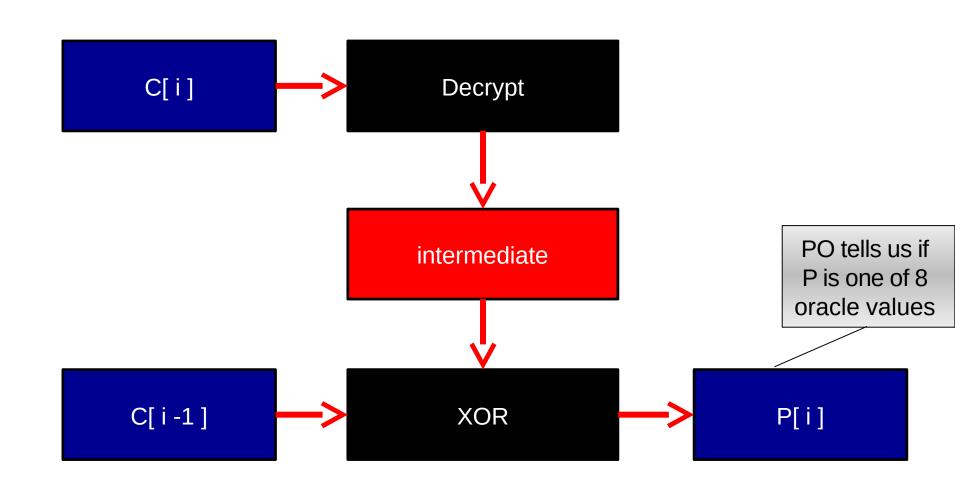
- Padding: terminating last plaintext block
 - -1337 => 1337060606060606 PKCS5Padding
 - -1337 => 133700000000000 ZeroPadding
- PKCS5 Padding Oracle
 - -"OK" if message ends with either of:
 - -0808080808080808 xx07070707070707
 - -xxxx060606060606 xxxxxx0505050505
 - -xxxxxxx04040404 xxxxxxxxx030303
 - -xxxxxxxxxxxx0202 xxxxxxxxxxxxxxx01



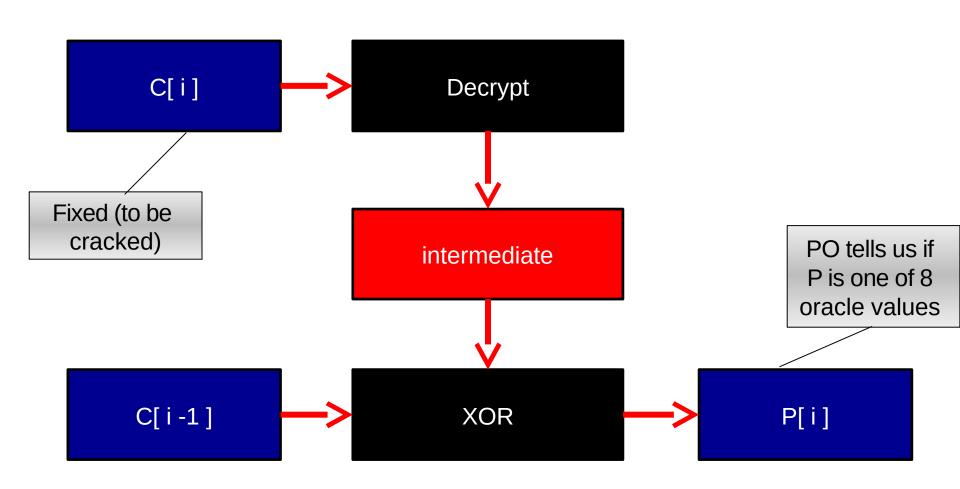




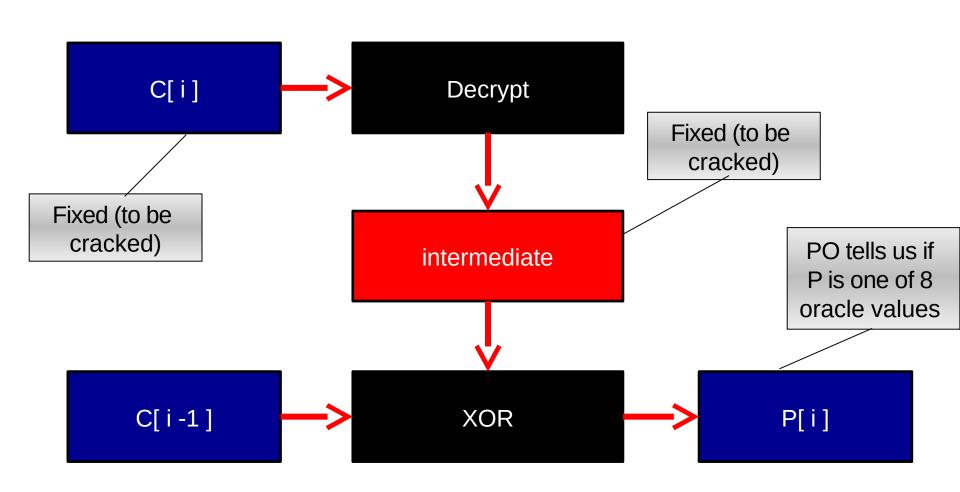




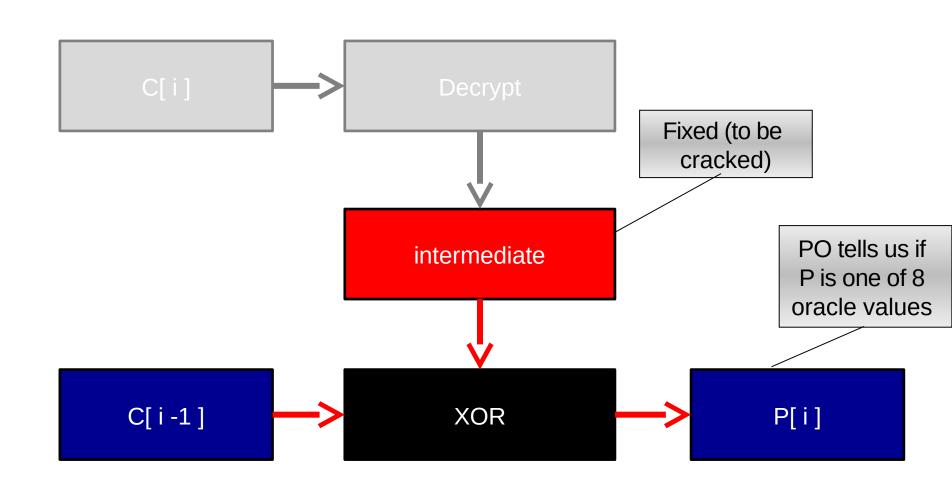




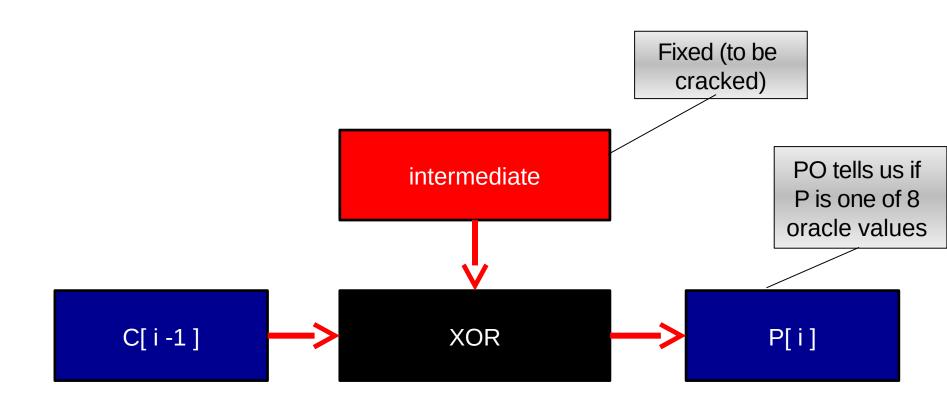




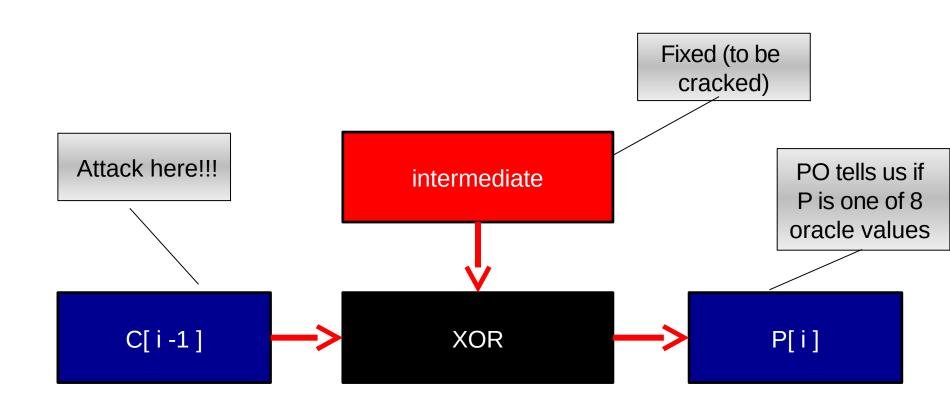












- CBC + PKCS5 Oracle Attack
 - Attacker replace C[i-1] with R, can choose any
 - -Attacker replace P[i] with unknown P
 - -intermediate = R XOR P when P=0808080808080808
 - -Padding Oracle: P is correct upon this R
 - -Padding Oracle leaks clues about intermediate



- CBC + PKCS5 Oracle Attack
 - -Start with R = 0 (or a random number), change LSB until Padding Oracle says OK!
 - $-\sim$ 128 messages find R: P = xx01 => OK!
 - $-\sim$ 256 messages find R: P = xx0202 => OK!
 - -~1024 messages find R: P = 0808080808080808 => OK!
 - -intermediate = R xor 0808080808080808
 - -P[i] = C[i] xor intermediate



- Why did padding Oracle work?
- SSL/TLS CBC cipher suits are broken
 - -Authenticates AFTER decryption
- Real long term solution: new cipher suits
 - –Authenticated Encryption (AEAD)
 - -Authenticate BEFORE decryption



- How was Padding Oracle fixed?
- Workaround: obfuscate the padding oracle
 - Decode as if Zero Padded
 - -Failed padding => slightly longer plaintext
 - -Calculate MAC
 - -Fail with nearly the same execution time
 - -pretends MAC failed, not padding



TLS RFC;

"This leaves a small timing channel, since MAC performance depends to some extent on the size of the data fragment, but it is not believed to be large enough to be exploitable, due to the large block size of existing MACs and the small size of the timing signal."



Lucky 13 (2013)

- Padding Oracle returns!
 - -SSL header size (13 bytes) is nearly optimal for attacking the ZeroPad + HMAC timing channel
- The small time difference can be detected
 - -Many attempts & Statistical models
 - Low latency between attacker and victim required
- Second Padding Oracle / Lucky13

Peter Magnusson, Twitter: @Blaufish_
Open Web Application
Sakerhetspodcasten.se & Omegapoint.se
Society Project Lex Coding solution to ensure constant

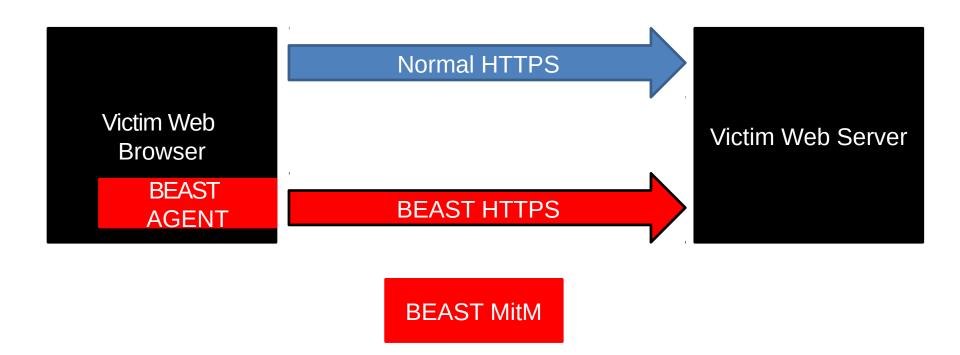
BEAST

Cryptanalytic attacks against SSL





BEAST (2011)



BEAST (2011)

- Browser Assisted Exploitation Against SSL/TLS
 - -Decrypts a message, usually a HTTP Cookie.
 - -Fools the browser to attack its own encryption
- Clever SOP rule bypass
 - -BEAST agent injected into HTTP://victim
 - -BEAST agent attacks HTTPS://victim
 - BEAST agent doesn't know cookie itself, but it is included in HTTPS sent from BEAST agent



BEAST (2011)

- Browser Assisted Exploitation Against SSL/TLS
 - -aka "Here come the XOR Ninjas".
- CBC Cipher Block Chaining mode
- Similar XOR-attack as seen in padding oracle
 - -Send P* = C* XOR IV XOR (R||i)
 - If one byte of $C^* == C$, we know 1 byte of real

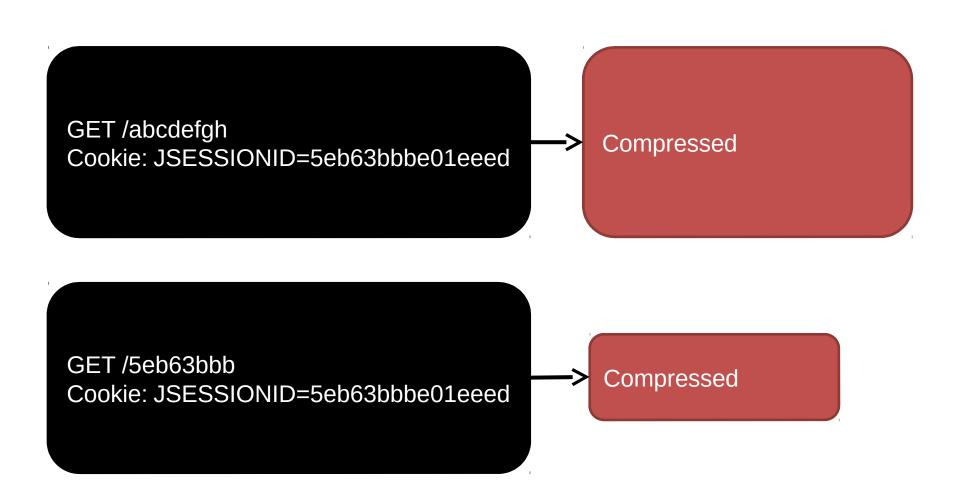
CRIME & BREACH

Cryptanalytic attacks against SSL





CRIME (2012)



CRIME (2012)

- Compression Ratio Info-leak Made Easy
- Attack HTTPS SSL/TLS compression
 - -GET/abc Cookie: JSESSIONID=5eb63bbbe01eeed
 - -GET/5eb Cookie: JSESSIONID=5eb63bbbe01eeed
 - -Shorter compressed message == guess is better
- SOP bypass not needed to launch CRIME

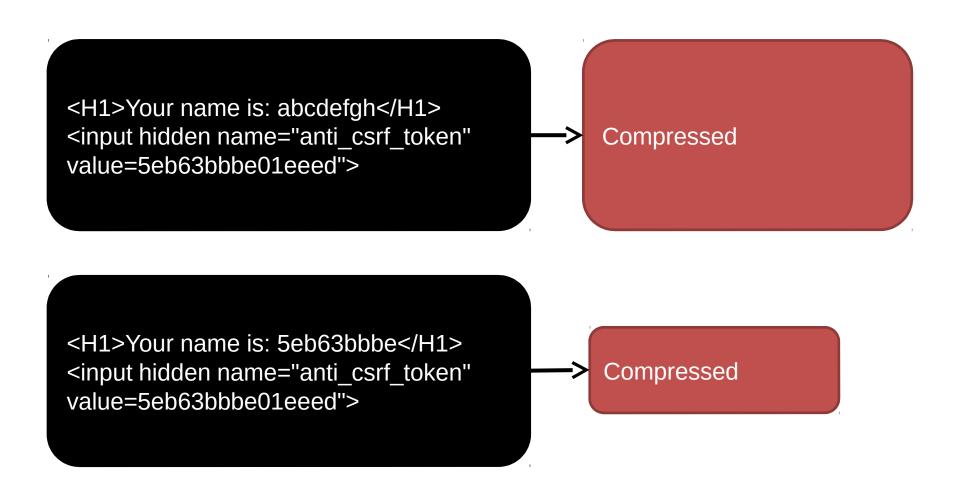
CRIME (2012)

- Various SSL changes to mitigate CRIME
 - -Many browsers refusing SSL Compression
 - -Some web servers refusing SSL Compression
- Jay! CRIME fixed, all done
 - But... SSL Compression is not the most common form of compression





BREACH (2013)



BREACH (2013)

- BREACH attacks HTTP response body:
 - -HTTP Compression must be enabled
 - -Response contains a secret
 - -Response reflects attacker input
- CRIME vs BREACH:
 - Both use compression oracles to decryptHTTPS
 - -CRIME exploit SSL/TLS compression of request
 - -BREACH exploit HTTP compression of



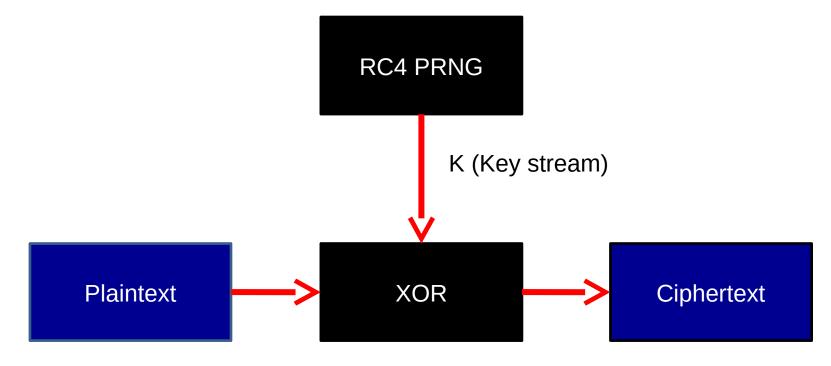
Weak Algorithms

Cryptanalytic attacks against SSL



RC4 (2013)

Designed 1987, very fast & useful





RC4 (2013)

- Known to be weak since 2001
 - -Bias (not uniform random) key stream
 - -Key stream leaks some of internal state
 - –WEP completely broken due to RC4 & bad design
- ...but RC4 flaws did not affect SSL/TLS (?)



RC4 (2013)

- Borderline practical attacks emerge
- RC4 for SSL/TLS can be broken
 - Assuming a lot of messages
 - -Assuming a lot of time
 - –ANY improvement on these attacks => practical
- There is no dirty workaround for RC4
 - -Not a SSL/TLS protocol issue
 - -Fundamental algorithm flaw, it must be removed





Agenda

What is HTTPS / SSL

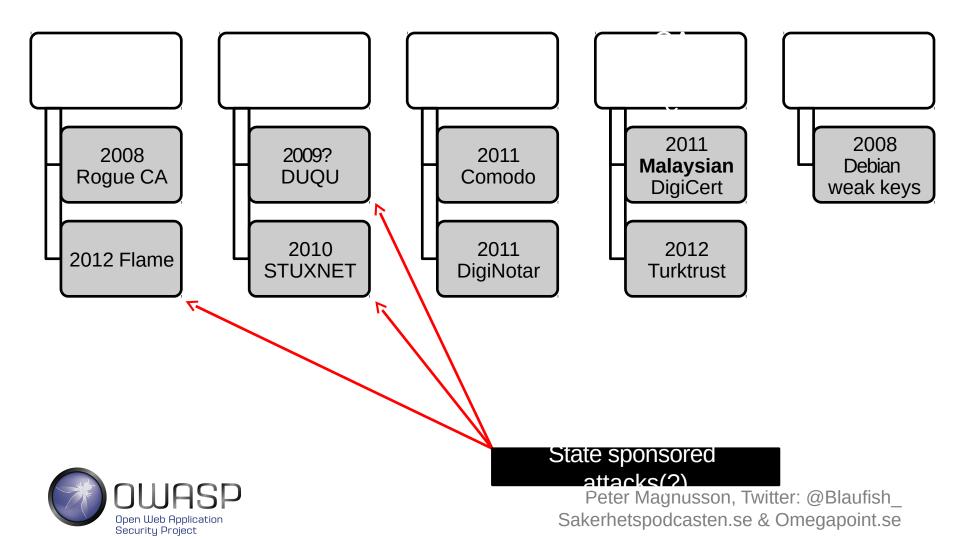
Attacks on SSL

Attacks on related technology

Coding SSL/TLS

Related attacks



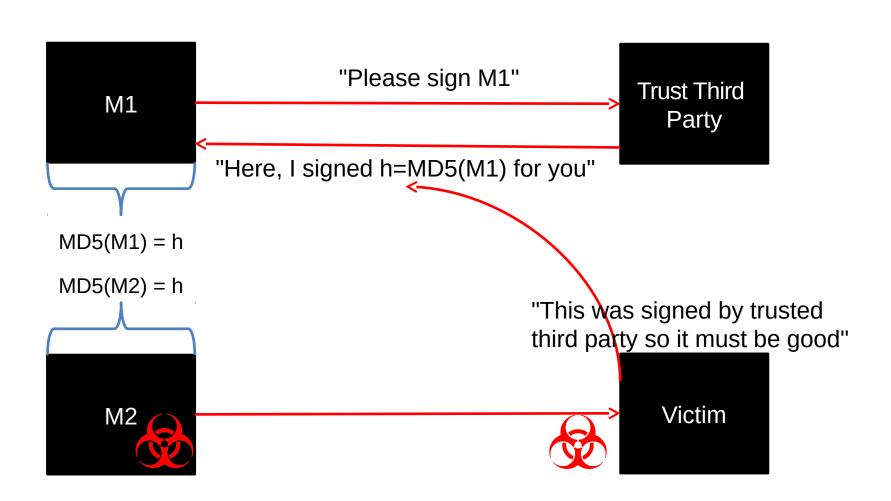


Related attacks: MD5 Preimage





MD5 Pre-image



Rogue CA certificate (2008)

- Academic team of security researchers
- Supercomputer: 200+ PlayStation 3
- MD5 second-preimage attack
 - -M1 != M2, H(M1) == H(M2)
 - -RapidSSL, MD5
 - -M1 = basic constraint CA=FALSE
 - -M2 = basic constraint CA=TRUE



Flame (2012)

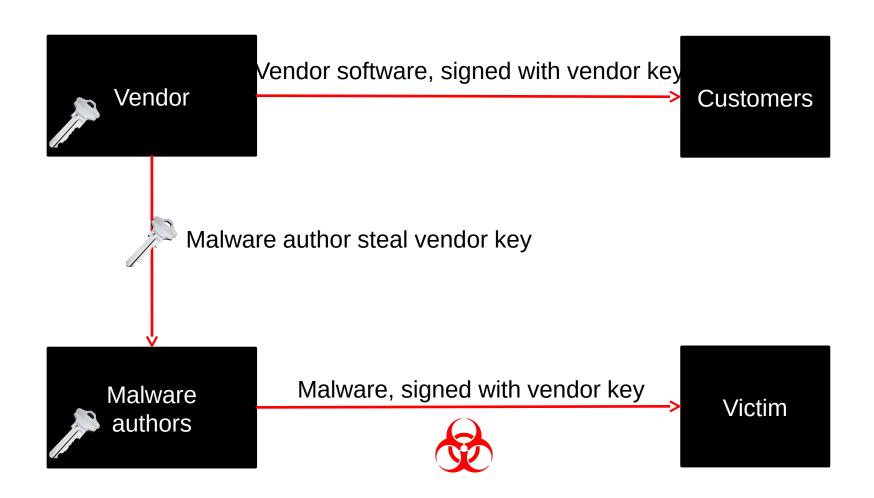
- State sponsored malware?
- MD5 second-preimage attack
 - -M1 != M2, H(M1) == H(M2)
 - -M1 = Microsoft Terminal Services license
 - -M2 = Code signing cert valid in Window Update
- Similar but NOT SAME as Rouge CA attack!
- Cryptanalyst team required to perform

Related attacks: Stolen keys





Stolen keys



DUQU – stolen crypto keys

- State sponsored malware?
- Signed with stolen code sign certificate
 - -C-Media Electronics, Inc. (certificate issued 2009, signature not time stamped)
- Purpose
 - –Key logger & attacks targeting small CA's according to McAfee. Espionage?



STUXNET – stolen crypto keys

- State sponsored malware?
- Signed with stolen code signing certificate
 - -Realtek Semiconductor Jan 25, 2010
 - -JMicron Technology Corp July 14, 2010
- Purpose
 - -Disrupt Iranian nuclear enrichment

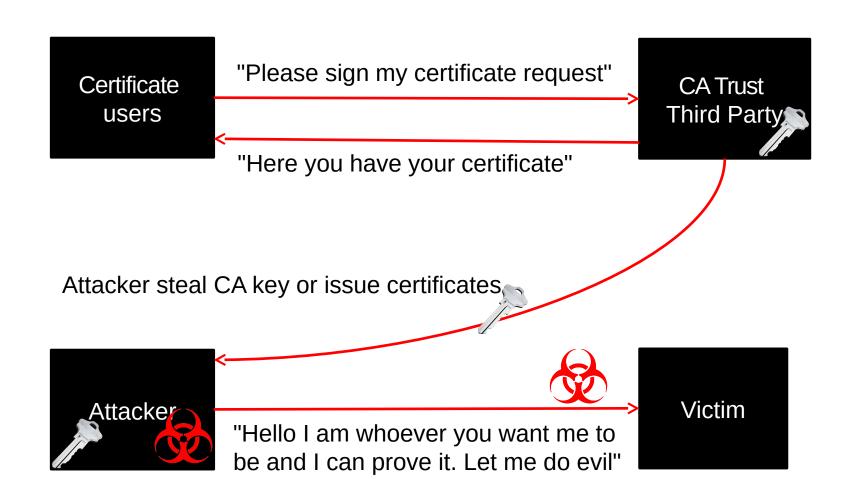


Related attacks: CA Breach





CA Breach



Comodo (2011)

- State sponsored attack, or only hackers?
- An affiliate Registration Authority (RA) hacked
- 9 certificates fraudulently issued
 - -Revoked



DigiNotar (2011)

- State sponsored attack or only hackers?
- Iranian Man in the Middle attack
 - MitM detected by Google Chrome browsers
 - -DigiNotar hacked by Iranian hackers
- 531 or more fake certificates signed
- Trust revoked, DigiNotar files for



Related attacks: CA Malpractice





CA Malpractice



Digicert Malaysia (2011)

- Not an attack; severe malpractice
 - -512 bit RSA keys (crackable by ordinary criminals)
 - –Certs issued w/o extensions (effectively CA, code signing, Server auth, etc... "do everything")
- Trust revoked for Digicert Malaysia
 - -Small CA name squatting on Digicert. Why would browsers etc allow different CAs with

samesame?
Peter Magnusson, Twitter: @Blaufish_
-Some zealots removed all Digicert, breaking these

Turktrust (2012 XMAS!)

- Not an attack; severe malpractice
 - –Accidentally issued intermediate CA to Turkish gov in 2011. Didn't clean up when becoming aware.
- Accidental Man-in-the-Middle Attack
 - -Christmas December 24 2012
 - -Turkish gov accidentally install intermediate CA from Turktrust into SSL inspecting firewall
 - -Google Chrome users alert Google to MitM





Trust Post Snowden



Snowden revelations

- Open algorithms are in general good (hard for NSA)
 - AES, Diffie-Hellman, Curve25519, Blake
 - Communication security in general works
- NSA targets implementations
 - Looks for weaknesses
 - Influence/strong arm of implementations
 - SW, HW and systems
- NSA targets std development
 - NIST, ISO, ETSI/SAGE, IEEE
 - IETF (SSL/TLS)
- NSA targets data at rest/in use
 - Cloud services Google, Appe, MS etc

Goal:

Find or create easier methods to gain access

Math works!

Track users and communication



What is suspect?

- Random generators
 - Intel True RNG (Bull Mountain) in Ivy Bridge, Haswell
 - NIST specified Dual_EC_DRBG

Suspect backdoor now confirmed

- Algorithms
 - NIST DES/3DES
 - NIST SHA-1
 - NIST, IEEE Elliptic Curves
 - SAGE 3G, 4G algorithms
 - China SMS4, ZUC
 - Russia GOST
- Protocols
 - SSL/TLS
- Implementations
 - HSMs
 - Closed Source libs, applications, systems
 - MS, Cisco, IBM etc

NSA huge patent owner for ECC. Big influence in std.

Can we trust security that is not open and has no transparent background?



What is the risk?

- Random generators
 - Weak keys, no real random numbers
- Algorithms
 - Weak algorithms not expected strength
- Protocols
 - Key leakage, session hijacking
- Implementations
 - Backdoors, unauthorized access

Is NSA really the main adversary?

Weaknesses are blind can be used by other adversaries



What is being done?

- Random generators
 - Push to open entropy source for Bull Mountain
 - Bull Mountain not replacing CSPRNG in Linux
 - Create open, transparent HSM
- Algorithms
 - Push to define new EC curves
 - Replace RC4 in SSL/TLS with modern, open stream cipher
 - Move away from dependency of NIST

NIST has lost a huge amount of trust

SHA-3

- Protocols
 - Reevaluation of SSL/TLS and other IETF sec standards
- Implementations
 - Several efforts to audit, evaluate, validate sec implementations
 - MS scramble to regain trust

Fundraiser to audit TrueCrypt
https://www.fundfill.com/fund/TrueCryptAudited
16k USD raised to date



What can you do?

- Random generators and keys
 - Test you generator and generated keys
 - Use longer keys (though it has costs)
- Algorithms
 - Move away from RC4, DES/3DES AES
 - Move away from MD5, SHA-1 to SHA-2 (256, 384, 512)
 - Be wary of ECC.

Perfect Forward Secrecy uses ECDSA

- Protocols
 - Move towards TLS 1.x
- Implementations
 - Use open implementations and libs
 - Test your applications and systems
 - Consider where you store (systems, services) stuff

Normal security strategy

Need to be done anyway



Agenda

What is HTTPS / SSL

Attacks on SSL

Attacks on related technology

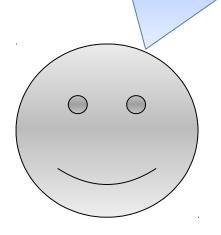
Coding SSL/TLS

Developing SSL/TLS code is hard

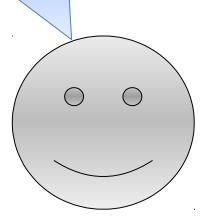


Developing SSL/TLS code is hard

Please make changes to this huge security stack which is poorly documented, you barely understood, and throws strange error messages upon failures



Sure! How hard can it be?



Why is developing SSL/TLS code hard?

- There are so many security aspects to authentication and SSL/TLS, that very few fully understand all of it.
 - -Authentication: There ~ 14 rules to be checked, probably more! http://en.wikipedia.org/wiki/Certification_path_validation_algorithm
 - -Failure: everything works great (only insecure)
- Your defaults are HOPEFULLY pretty



Why is developing SSL/TLS code hard?

- BE PARANOID
 - -Test test test
 - -Think through many times what the code does
- Understand incomprehensive errors
 - -don't rush code changes to deal with SSL errors
 - -SSL error are usually logical, technical and security oriented. Unhelpful but CAN be understood.



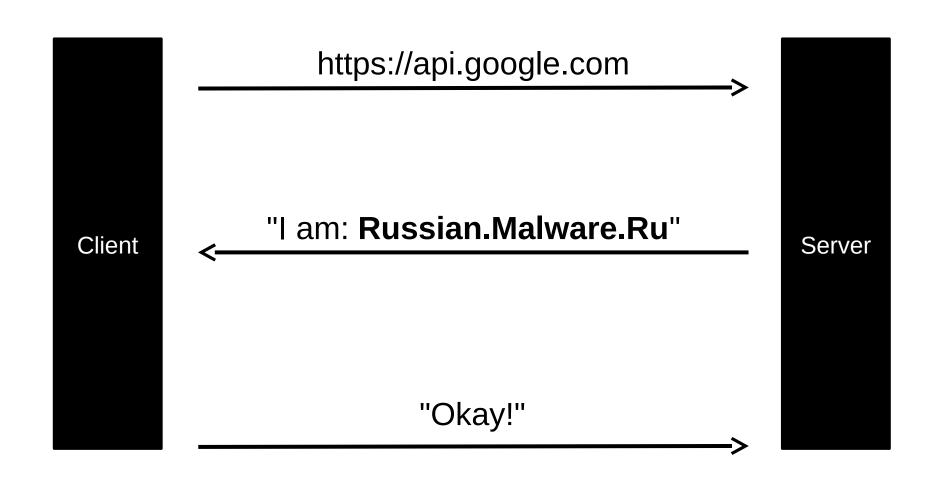
Classic WTF Horror

BEWARE this section contain DO NOTs you should NOT copy into your own code





Classic WTF Horror



Classic WTF Horror

- Any valid certificate for ANY site may spoof the site
- Sadly, this is "best answer" in various



Stacks and libraries





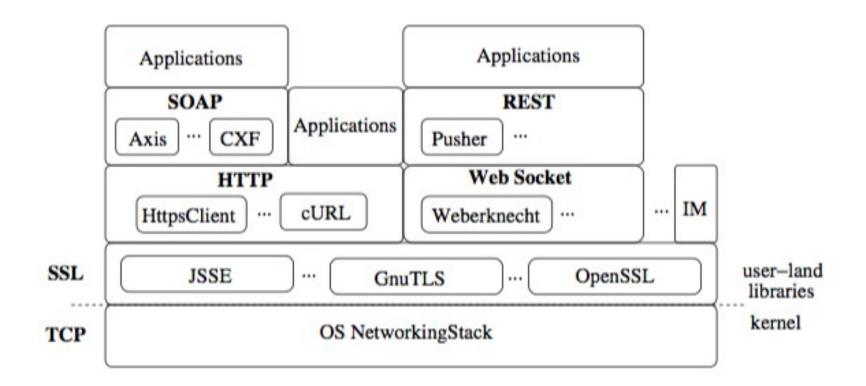


Figure 2: Protocol stack.

SSL used through layers with different APIs



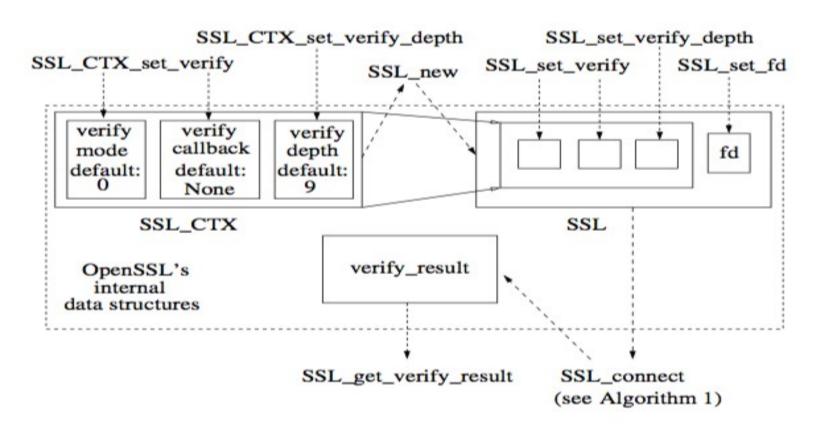


Figure 3: OpenSSL API for setting up SSL connections with the default chain-of-trust verification.

Complex APIs makes for easy mistakes



This interface is almost perversely bad. The VERIFYPEER parameter is a boolean, while a similar-looking VERIFYHOST parameter is an integer. The following quote from the cURL manual explains the meaning of CURLOPT_SSL_VERIFYHOST:

1 to check the existence of a common name in the SSL peer certificate. 2 to check the existence of a common name and also verify that it matches the hostname provided. In production environments the value of this option should be kept at 2 (default value).

Well-intentioned developers not only routinely misunderstand these parameters, but often set CURLOPT_SSL_VERIFY HOST to TRUE, thereby changing it to 1 and thus accidentally disabling hostname verification with disastrous consequences (see Section 7.1).

Validation control in cURL

```
curl_setopt($curlHandle, CURLOPT_SSL_VERIFYPEER, true);
curl_setopt($curlHandle, CURLOPT_SSL_VERIFYHOST, true);
```

. . .

Example code from Amazon FPS (PHP)



7.3 PayPal IPN in ZenCart

ZenCart's functionality for PayPal IPN shows a profound misunderstanding of cURL's parameters. It disables certificate validation entirely, yet attempts to enable hostname verification—even though the latter has no effect if certificate validation is disabled.

Even worse example code provided by PayPal



```
def _verify_hostname(self, hostname, cert):
  # Verify hostname against peer cert
  # Check both commonName and entries in subjectAltName,
 # using a rudimentary glob to dns regex check
  # to find matches
  common_name = self._get_common_name(cert)
  alt_names = self._get_subject_alt_names(cert)
  # replace * with alphanumeric and dash
  # replace . with literal .
 valid_patterns = [re.compile(pattern.replace(r".", r"
      \.").replace(r"*", r"[0-9A-Za-z]+"))
   for pattern
      in (set(common name) | set(alt names))
  return any (
   pattern.search(hostname)
   for pattern in valid_patterns
```

Hostname validation in Apache Libcloud

google.com = *oogle.com moogle.com, scroogle.com are **BAD** domains

Handling SSL/TLS dev challenges



Handling SSL/TLS dev challenges

- Test environments
 - –Don't create code to "handle SSL errors for dev/test"
 - -Start dev environments with a Test trust store which trust the test environments
 - -Run with secure standard code in dev/test
- Special cases
 - If you absolutely have to mess around, mess with specific connections instead of changing



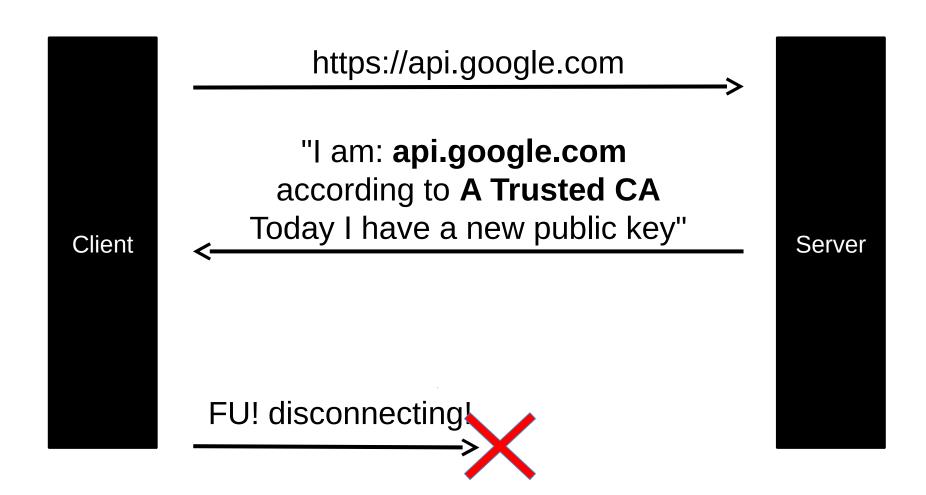
Cerificate Pinning

https://www.owasp.org/index.php/ Certificate_and_Public_Key_Pinni ng#Android





Certificate Pinning



Certificate Pinning

- Force your system to only trust a specific certificate
 - -If you do not trust certificate authorities
 - -CA breaches
 - -CA malpractice
 - Compelled certificates by a malicious government
- But only add security, don't remove checks
 - -https://www.owasp.org/index.php/Certificate_an

Certificate Pinning (1/2)

```
private static String PUB_KEY =
"30820122300d06092a864886f70d0101 .. .. .. .. ...":
public void checkServerTrusted(X509Certificate[] chain, String authType)
throws CertificateException {
  // Perform customary SSL/TLS checks
  try {
   TrustManagerFactory tmf = TrustManagerFactory.getInstance("X509");
   tmf.init((KeyStore) null);
   for (TrustManager trustManager : tmf.getTrustManagers()) {
    ((X509TrustManager) trustManager).checkServerTrusted(chain,
authType);
```



Certificate Pinning (2/2)

```
// Hack ahead: BigInteger and toString(). We know a DER
encoded Public Key begins
// with 0x30 (ASN.1 SEQUENCE and CONSTRUCTED), so
there is no leading 0x00 to drop.
    RSAPublicKey pubkey = (RSAPublicKey)
chain[0].getPublicKey();
```

String encoded = new BigInteger(1 /* positive */, pubkey.getEncoded()).toString(16);

Force crazy TLS Security?



Force crazy TLS Security?

SecureSocketFactory extends **SSLSocketFactory** {
public Socket **createSocket**(Socket s,) throws
IOException {

s.setEnabledProtocols(new String[] { "TLSv1.2" });

AES GCM and other AEAD cipher suites not officially supported in java7





SSL Best Practice



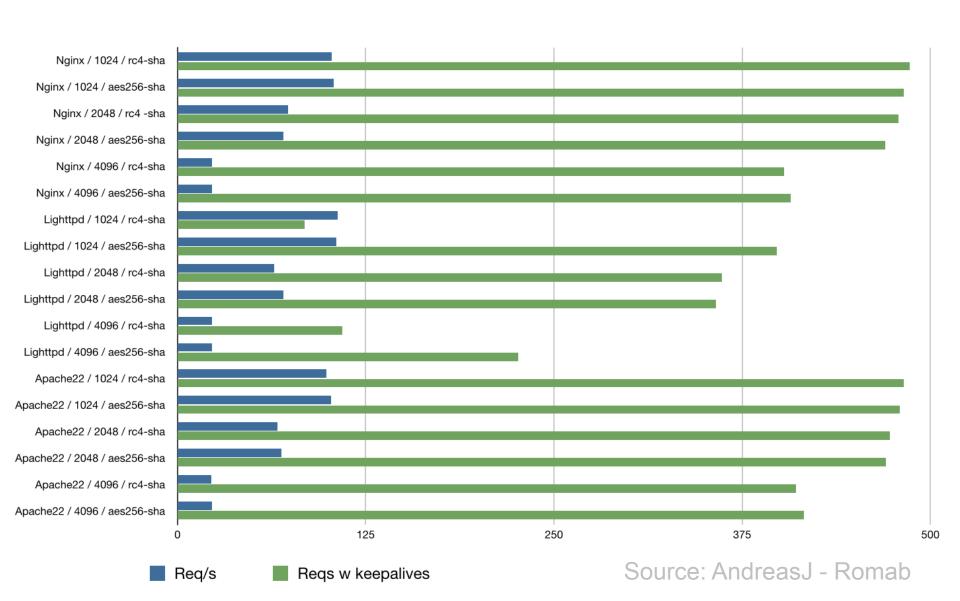
USE HTTPS!



Cost of longer keys

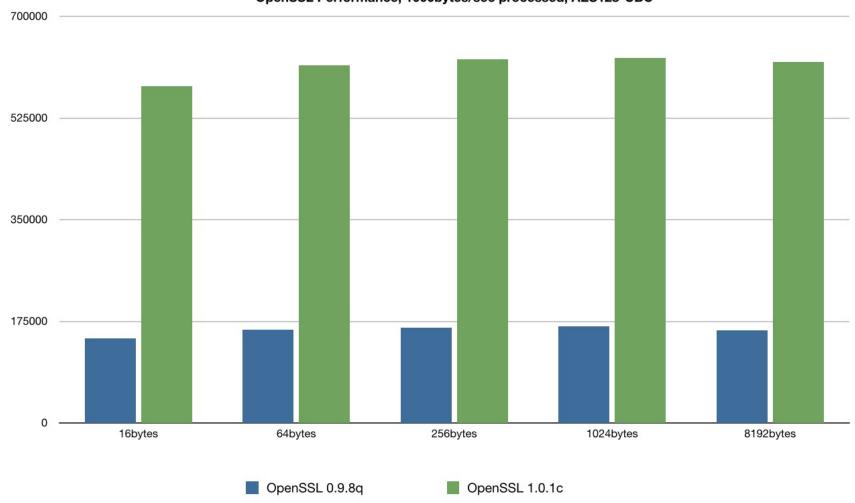
- Rapidly decreased performance
 - But Keepalives is more important
 - Upgrade of openssl can improve performance
- Interoperability
 - OSX, iOS supports max 4096 bits by default
 - Chrome, Android max ~2300 bits until recently

server/rsa key/cipher



OpenSSL

OpenSSL Performance, 1000bytes/sec processed, AES128-CBC

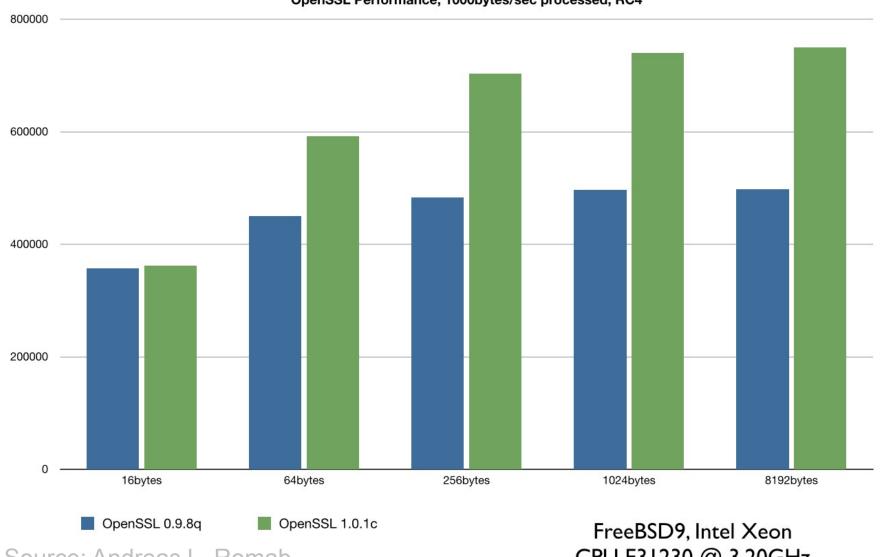


FreeBSD9, Intel Xeon CPU E31230 @ 3.20GHz

Source: AndreasJ - Romab

OpenSSL

OpenSSL Performance, 1000bytes/sec processed, RC4



Source: AndreasJ - Romab

CPU E31230 @ 3.20GHz



CA, Trust Store requirements

Microsoft

Microsoft Security Advisory (2880823)

Deprecation of SHA-1 Hashing Algorithm for Microsoft Root Certificate Program

Published: Tuesday, November 12, 2013

Version: 1.0

General Information

Executive Summary

Microsoft is announcing a policy change to the Microsoft Root Certificate Program. The new policy will no longer allow root certificate authorities to issue X.509 certificates using the SHA-1 hashing algorithm for the purposes of SSL and code signing after January 1, 2016. Using the SHA-1 hashing algorithm in digital certificates could allow an attacker to spoof content, perform phishing attacks, or perform man-in-the-middle attacks.

Recommendation: Microsoft recommends that certificate authorities no longer sign newly generated certificates using the SHA-1 hashing algorithm and begin migrating to SHA-2. Microsoft also recommends that customers replace their SHA-1 certificates with SHA-2 certificates at the earliest opportunity. Please see the **Suggested Actions** section of this advisory for more information.

No root certs **issued** with SHA-1 after 2016

Use certs with SHA-2 (there are broblems though)



CA, Trust Store requirements

Mozilla

- June 30, 2011 Mozilla will stop accepting MD5 as a hash algorithm for intermediate and end-entity certificates. After this date software published by Mozilla will return an error when a certificate with an MD5-based signature is used.
 - bug 650355 "Stop accepting MD5 as a hash algorithm in signatures (toggle security.enable_md5_signatures to false)" -- Fixed in Mozilla 16 (Firefox 16).
 - bug 590364 "By default, stop accepting MD5 as a hash algorithm in certificate signatures" Until this bug is fixed, non-Gecko software that uses NSS will still accept MD5 signatures. (Gecko is the layout engine developed by the Mozilla Project, originally called NGLayout.) -- In NSS 3.14
- December 31, 2013 Soon after this date, Mozilla will disable the SSL and Code Signing trust bits for root certificates with RSA key sizes smaller than 2048 bits. If those root certificates are no longer needed for S/MIME, then Mozilla will remove them from NSS.
 - TEST: You can test the behavior of changing the root certificate trust bit settings, as described here: https://wiki.mozilla.org/ /CA:UserCertDB#Changing Root Certificate Trust Bit Settings

No root certs with 1024 bit keys in January 2014

Check your root



Deploying Forward Secrecy

- Upgrade server and libs
 - Apache 2.4
 - Nginx 1.0.6, 1.1.0
 - OpenSSL 1.0.1c
 - GnuTLS 3.2.7

- Performance cost
 - Est 15% CPU and comm during init



TEST HTTPS!



Server validation



Home Qualys.com Pr

Projects

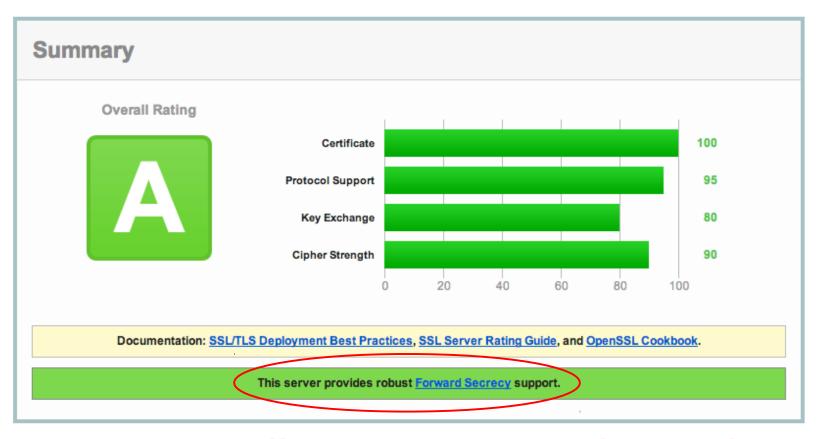
Contact

You are here: Home > Projects > SSL Server Test > romab.com

SSL Report: romab.com (192.195.142.8)

Assessed on: Thu Nov 21 13:11:35 UTC 2013 | Clear cache

Scan Another »



https://www.ssllabs.com/ssltest/





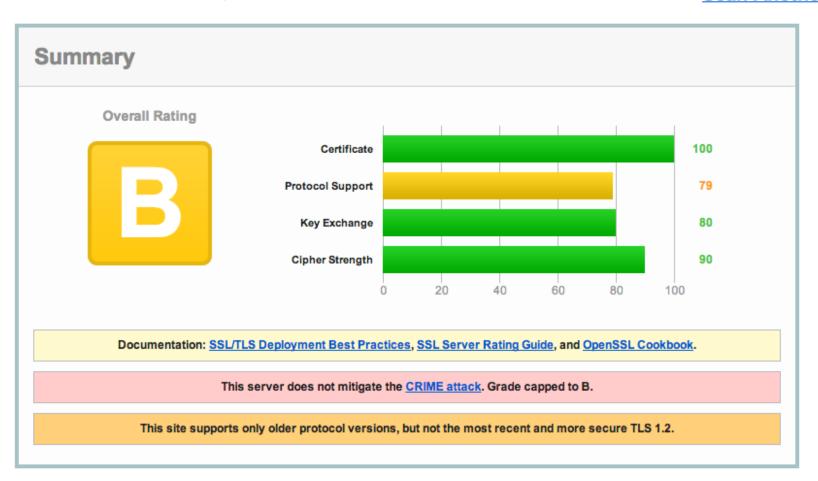
Home Qualys.com Projects Contact

You are here: Home > Projects > SSL Server Test > moIndal.se

SSL Report: molndal.se (31.193.200.10)

Assessed on: Thu Nov 21 14:22:19 UTC 2013 | Clear cache

Scan Another »







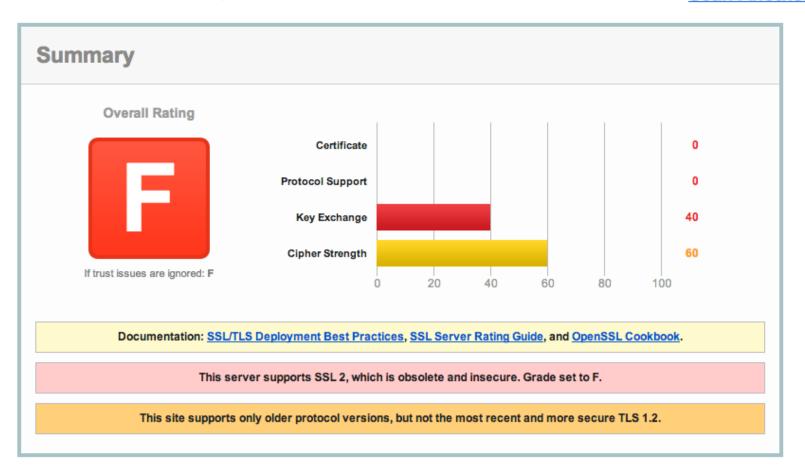
Home Qualys.com Projects Contact

You are here: Home > Projects > SSL Server Test > abanan.se

SSL Report: abanan.se (62.101.37.2)

Assessed on: Thu Nov 21 14:19:16 UTC 2013 | Clear cache

Scan Another »







Home Qualys.com Projects Contact

You are here: Home > Projects > SSL Server Test > omegapoint.se

SSL Report: omegapoint.se (194.103.93.60)

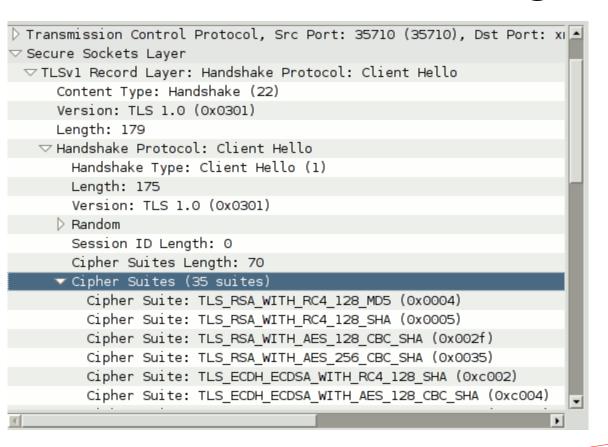
Assessed on: Thu Nov 21 14:04:12 UTC 2013 | Clear cache

Scan Another »

Assessment failed: No secure protocols supported



OWASP Check & change in-app suites



Android uses the cipher suite order in Java

AES-256 and SHA1 to RC4 and MD5 in 2010

Android 2.2.1	Android 2.3.4, 2.3.7	Android 4.2.2, 4.3
DHE-RSA-AES256-SHA	RC4-MD5	RC4-MD5
DHE-DSS-AES256-SHA	RC4-SHA	RC4-SHA
AES256-SHA	AES128-SHA	AES128-SHA
EDH-RSA-DES-CBC3-SHA	DHE-RSA-AES128-SHA	AES256-SHA

Change the order in your app

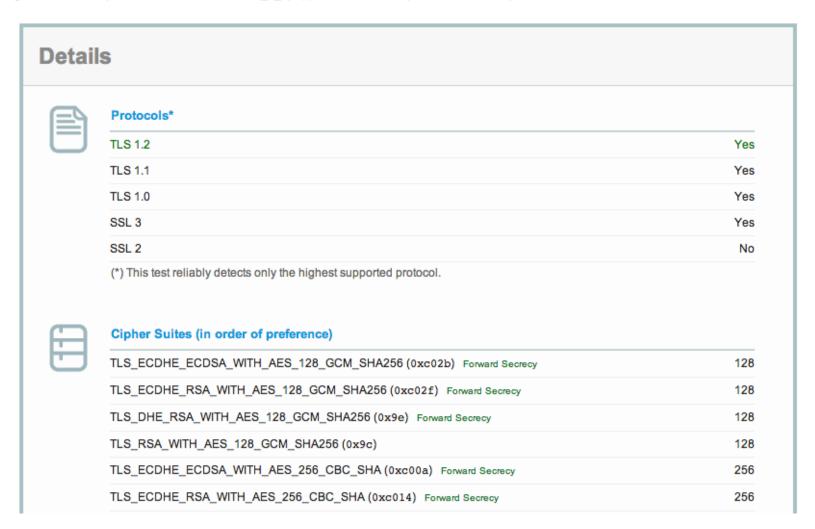
http://op-co.de/blog/posts/android ssl downgrade/



Client validation

SSL/TLS Capabilities of Your Browser (Experimental)

User Agent: Mozilla/5.0 (Macintosh; Intel Mac OS X 10_9_0) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/31.0.1650.26 Safari/537.36



https://www.ssllabs.com/ssltest/viewMyClient.html



Client Stress Testing

- TLSpretense
 - Check that your cliedbes not fail open
 - Accepts certs for wrong domain
 - Accepts broken chains
 - Null byte host name



Sslyze

- Open source SSL/TLS testing tool
 - Test public AND private servers
- IIS sponsored development of new functions
 - OCSP, CRL, Multiple Trust Stores, SNI etc
 - HSTS, HTTP vs HTTPs (content)
- Upstream merge of some parts
 - New generation with API breakage during project



Sslyze

>./sslyze.py --regular --crl --ocsp --sni=auto --hsts --starttls=auto www.kirei.se

* Certificate: CRL verification: Certificate not revoked in CRL SNI: SNI enabled with virtual domain www.kirei.se Trusted or NOT Trusted: Trusted OCSP verification: Certificate not revoked Validated by Trust Store: apple-20130529 Validated by Trust Store: microsoft-20130905 Validated by Trust Store: copied_ca Validated by Trust Store: stripped_ca Validated by Trust Store: java7v25-20130810 Validated by Trust Store: mozilla-20130812 X509 Policy in certificate: True Hostname Validation: OK - SNI CN www.kirei.se Matches SHA1 Fingerprint: FB35A7AAEC8A32FCC7E4016EAF5734459DCF5809 Common Name: www.kirei.se /C=GB/ST=Greater Manchester/L=Salford/O=Comodo CA Limited/CN=PositiveSSL CA Issuer: Serial Number: BC95078646835C5C090AAB839F1DDBFE Not Before: Aug 25 00:00:00 2011 GMT Not After: Aug 24 23:59:59 2014 GMT Signature Algorithm: sha1WithRSAEncryption Key Size: 2048 X509v3 Subject Alternative Name: DNS:www.kirei.se, DNS:kirei.se

* HSTS :

Supported:

max-age=864000



./sslyze.py --regular --sni=auto --starttls=auto --crl --ocsp --hsts ab.se

Certificate: No CRL URI in certificate CRL verification: SNI enabled with virtual domain ab.se SNI: Trusted or NOT Trusted: NOT Trusted OCSP verification: No OCSP Responder in certificate Not validated by Trust Store: apple-20130529 - self signed certificate Not validated by Trust Store: microsoft-20130905 - self signed certificate Not validated by Trust Store: copied_ca - self signed certificate Not validated by Trust Store: stripped_ca - self signed certificate java7v25-20130810 - self signed certificate Not validated by Trust Store: mozilla-20130812 - self signed certificate Not validated by Trust Store: X509 Policy in certificate: False Hostname Validation: MISMATCH SHA1 Fingerprint: 28953D02EE7763D865242C2BE08E2F8DDDB610EE Common Name: intranet.ab.se Issuer: /CN=intranet.ab.se Serial Number: B365F5B94A020F50 Not Before: Jul 23 17:30:56 2010 GMT Not After: Jul 20 17:30:56 2020 GMT Signature Algorithm: sha1WithRSAEncryption Key Size: 1024



SSL Status

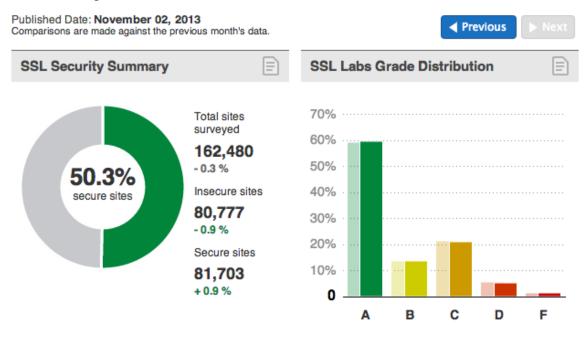
November 02, 2013



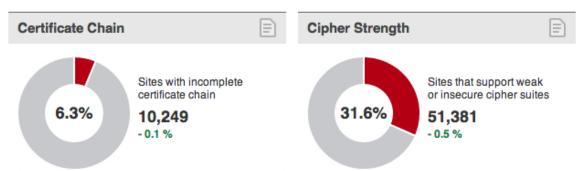
SSL Pulse 4

Survey of the SSL Implementation of the Most Popular Web Sites

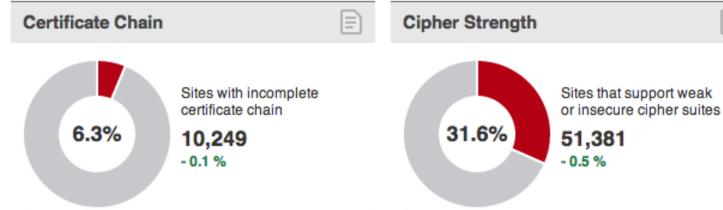
Summary

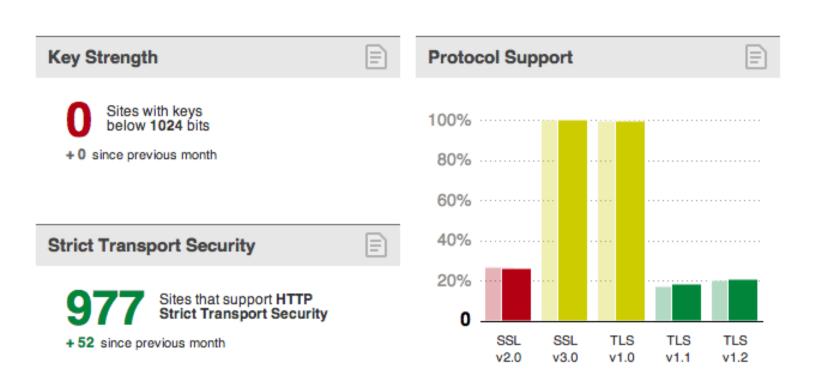


Key Findings









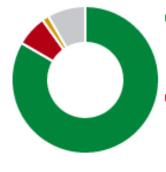
: https://www.trustworthvinternet.org/ssl-pulse/

Renegotiation Support



Key Strength Distribution





Secure renegotiation

135,533 83.4%

+ 0.4 %

Insecure renegotiation

11,020 6.8%

- 0.3 %

Both

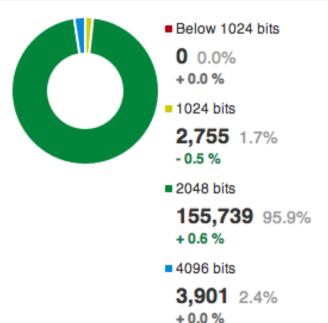
2,323 1.4%

+ 0.1 %

No support

13,604 8.4%

- 0.1 %



Extended Validation Certificates



Sites that have an Extended Validation certificate 14,743

+0.1%

BEAST Attack







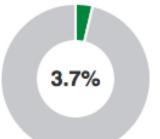
Peter Magnusson, Twitter: @Blaufish_ Sakerhetspodcasten.se & Omegapoint.se

SPDY



TLS Compression / CRIME





Sites that support the SPDY protocol

6,023 +0.7%



Forward Secrecy











Not supported

87,565 53.9%

- 0.1 %

Some FS suites enabled

67,847 41.8%

+0.0%

Used with modern browsers

6,016 3.7%

+ 0.1 %

Used with most browsers

1,052 0.6%

RC4



Not Supported

11,857 7.3%

+ 0.1 %

Some RC4 suites enabled

91,358 56.2%

- 0.1 %

Used with modern browsers

59,265 36.5%

+0.0%

+0.0%

Questions?

These slides will be available at:
www.slideshare.net/blaufish
www.owasp.org/index.php/Gothenburg





References

- OpenSSL Cookbook
 - https://www.feistyduck.com/books/openssl-cookbook/
- SSL/TLS Deployment Best Practice
 - https://www.ssllabs.com/projects/best-practices/index.html
- Sslyze
 - IIS version: https://github.com/kirei/sslyze
 - Upstream: https://github.com/iSECPartners/sslyze

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